

Agile Earth Observation Satellite Scheduling with a Quantum Annealer

Presentation at COMET SIL

DEFENCE AND SPACE

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AIRBUS

AGENDA

- ❑ Context
 - ❑ Trends for agile Earth Observation satellite scheduling
 - ❑ Promises of quantum computing based optimization
- ❑ EO satellite scheduling with a Quantum Annealer
- ❑ Benchmarking methodology and results
- ❑ Conclusion and perspectives

Trends for Earth Observation Mission Planning

Trends Very High Resolution Agile EO Satellites

- Constellation of satellites: from 2-4 to 10s to 100s of platforms
- Smaller instrument footprint → larger volume of candidate meshes (i.e. surface elements) to plan per programming period
- Enhanced agility: multiplication of acquisition opportunities and planning solutions
- Multi-Objective optimization: priority satisfaction, capacity (surface) maximization, age of information, weather conditions...



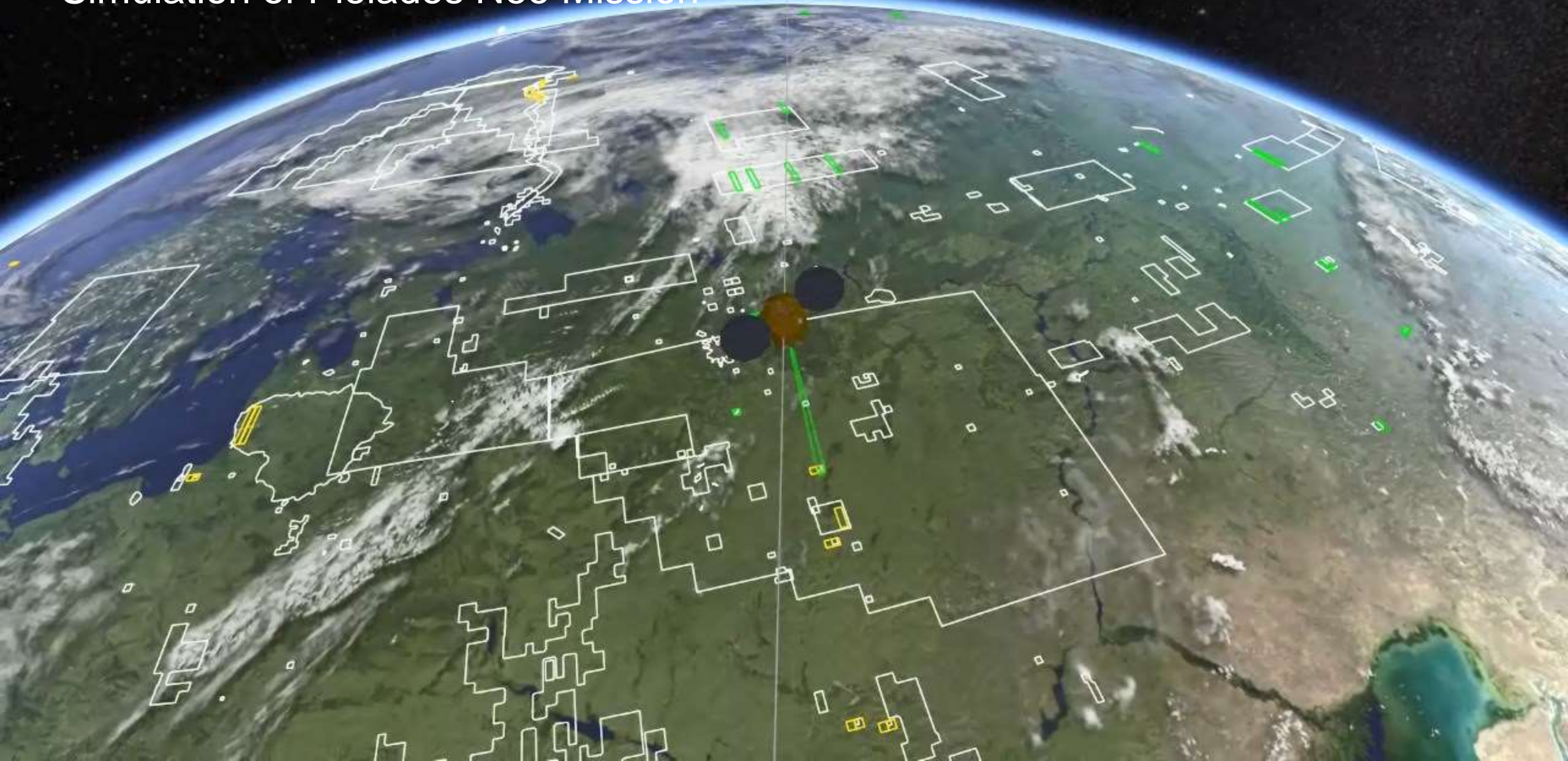
Bottom line

- EO Mission Planning is a well-known **multi-objective NP-hard optimization** problem **under uncertainty**
- Current trends indicate a **combinatorial explosion (# decision variables, # constraints)** for future Earth Observation systems

Expectations

- Current Mission Planning solutions are based on (sub-optimal) heuristic algorithms (greedy or dynamic programming)
- Experiments on smaller instances of the problem have shown gains ranging from 10% to 20% between the optimum and the solution obtained by current approximate algorithms

Simulation of Pleiades Neo Mission



EO Mission Planning Problem Statement

Mission Planning: Must determine an optimal acquisition plan for an Earth Observation satellite

Input data

- R is the set of acquisition requests, I_r is the set of imaging attempts for request $r \in R$.
- $\forall r \in R, \forall i \in I_r, w_{r,i}$ is the score of the imaging attempt
- $\forall r \in R, \forall i \in I_r, t_i^r$ is the start time of the imaging attempt

Decision variable

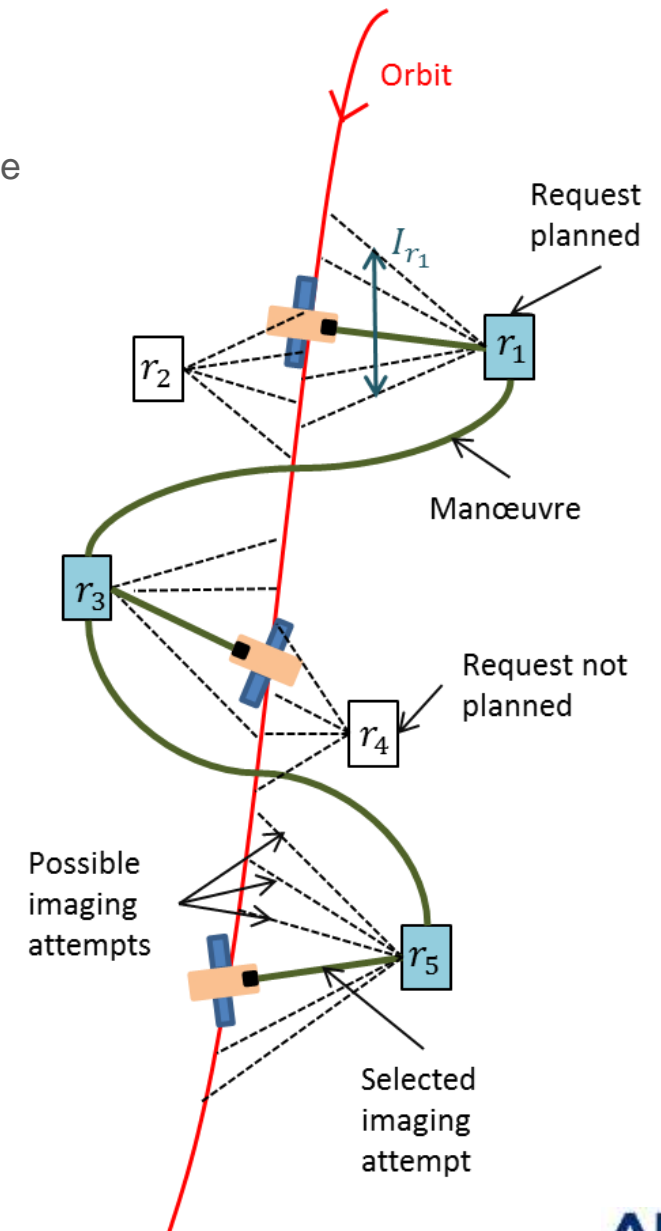
- $x_{r,i}$ is the binary variable indicating whether the candidate attempt i is selected in the plan
 - The number of binary variables is $N_{\text{variable}} = \sum_r |I_r|$

Constraints

- Maximally one assigned attempt i per request r : $\forall r \in R, \sum_{i \in I_r} x_{r,i} \leq 1$
- Some consecutive imaging attempts are not possible:
 - $F_{r_1, r_2} = \{(i, j) \in (I_{r_1}, I_{r_2}) \mid t_i^{r_1} \leq t_j^{r_2} \ \&\& \ t_j^{r_2} < t_i^{r_1} + T_i^{r_1, \text{acquisition}} + T_{i \rightarrow j}^{r_1, r_2, \text{maneuver}}\}$
 - $\forall (r_1, r_2) \in R^2, \text{with } r_1 \neq r_2, \forall (i, j) \in F_{r_1, r_2}: x_{r_1, i} \cdot x_{r_2, j} = 0$

Objective: Total score of the schedule

- Minimize $C = -\sum_r \sum_{i \in I_r} w_{r,i} x_{r,i}$

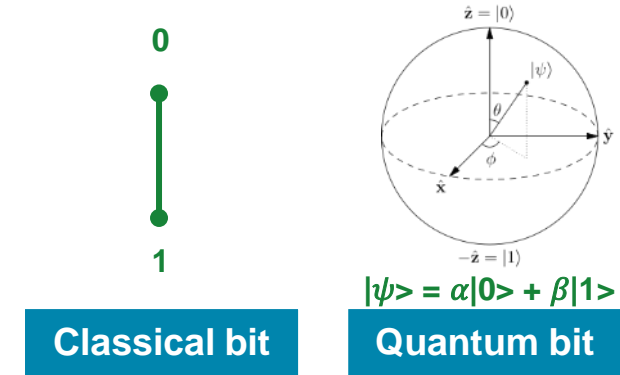


Quantum Computing in a Nutshell

Superposition Principle

- A qubit can be seen as a superposition of two basis vectors

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$
- A n-qubit register represents a 2^n -dimensional vector space, allowing for exponentially greater information processing



Classical vs Quantum Computing

	Steps	Classical Computation	Quantum Computation
1	Initial State	A single length- N binary string e.g. 00	Superposition of length- N binary strings e.g. $\frac{ 00\rangle + 11\rangle}{\sqrt{2}}$
2	Apply Operations	Logic gates e.g. 00 \rightarrow 01	Quantum logic gates e.g. $\frac{ 00\rangle + 11\rangle}{\sqrt{2}} \rightarrow \frac{ 10\rangle + 01\rangle}{\sqrt{2}}$
3	Read Out	Read out e.g. READ(01) \rightarrow OBSERVED: 01	Measure e.g. MEASURE($\frac{ 10\rangle + 01\rangle}{\sqrt{2}}$) \rightarrow OBSERVED: either 10 or 01

$2^n \times 2^n$ Unitary operators

$$\sum_i |x_i|^2 = 1$$

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1m} \\ a_{21} & a_{22} & \dots & a_{2m} \\ \vdots & \vdots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nm} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} d_1 \\ d_2 \\ \vdots \\ d_n \end{bmatrix}$$

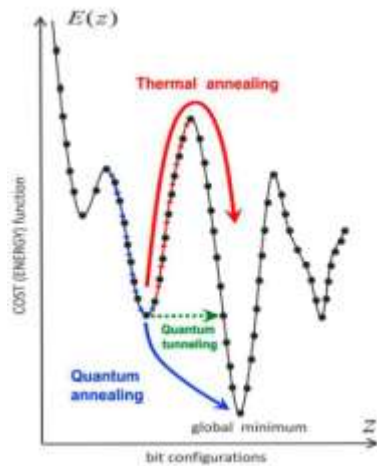
of components grows exponentially with # of qubits

unitary

Quantum Computing in a Nutshell

Quantum Annealing Computer (D-Wave)

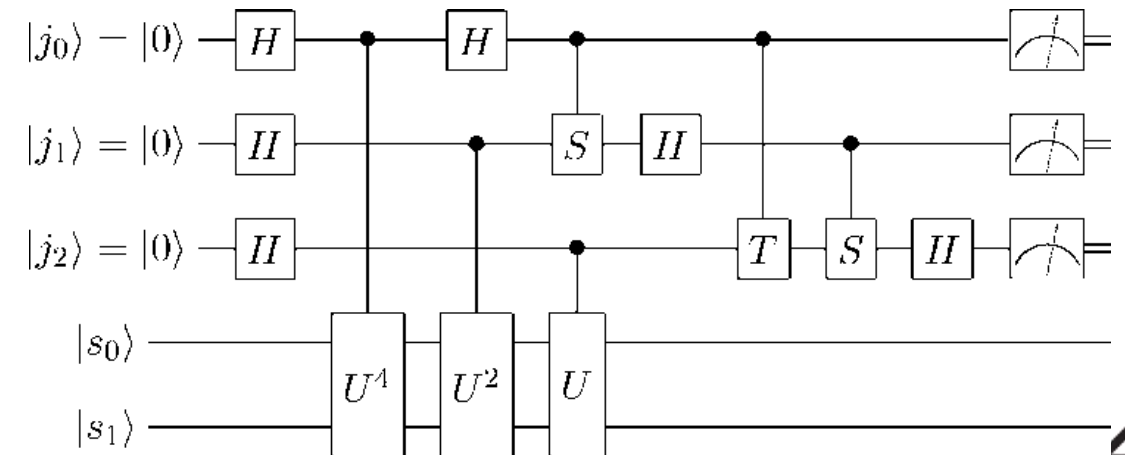
- Not a general purpose quantum computer, but uses quantum properties to solve discrete optimization problems
- Natural evolution of quantum-mechanical system (using quantum tunnelling) towards a ground state minimizing its energy



Quantum Annealer

General Purpose Quantum Computer (IBM, Google)

- Quantum circuits are composed of elementary gates and operate on qubits
- QC equivalent to classical boolean feed-forward networks, except they are reversible (i.e. quantum circuits can be evaluated in both directions)



Quantum Circuit

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- Promises of quantum computing based optimization

□ EO satellite scheduling with a Quantum Annealer

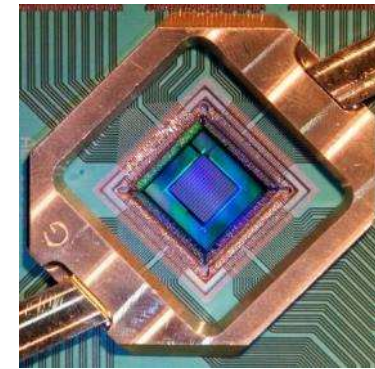
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Quantum Annealer in a Nutshell

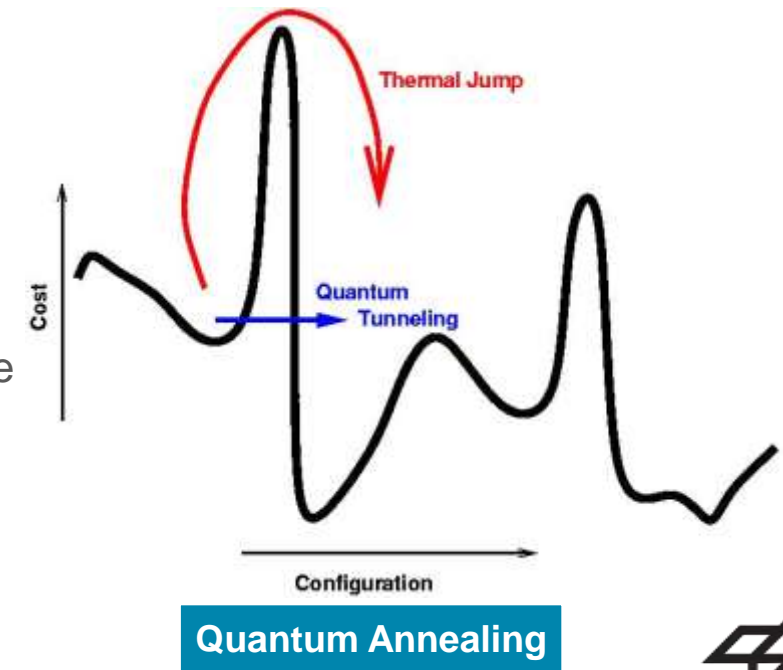
Quantum Computing for Combinatorial Optimization

Quantum Annealing

- Solves Quadratic Unconstrained Binary Optimization (QUBO) problems i.e. minimize $H(\{x_i\}) = \sum_i Q_{ii}x_i + \sum_{i<j} Q_{ij}x_ix_j$, where $x \in \{0, 1\}^N$
- Requires to formulate your discrete optimization problem as a QUBO
- QA can be seen as a stochastic process: several annealing runs are performed from a given initial state (e.g. uniformly distributed quantum superposition of all possible states)
- After a fixed elapsed time, the final state is measured providing a solution sample
- After a fixed number of runs, the solution sample having minimum energy is kept
- QA remains an approximate optimization technique, but the number of runs can be increased to reach a given probability of finding the exact solution



© D-Wave



EO Mission Planning Problem as a QUBO

Quadratic Unconstrained Binary Optimization (QUBO) formulation

- QUBO: $\min q(x) = x^T Q x = \sum_{j=1}^n Q_{j,j} x_j + \sum_{j,k=1}^n Q_{j,k} x_j x_k$ with Q an upper-triangular quadratic matrix
- Constraint equations in a quadratic form:
 - (1) : Max one imaging attempt per request : $C_u = \sum_r \sum_{i,j \in I_r, i < j} \min\{w_{r,i}, w_{r,j}\} x_{r,i} x_{r,j}$
 - (2) : Non feasible maneuver $C_t = \sum_{r_1, r_2} \sum_{i,j \in F_{r_1, r_2}} \min\{w_{r_1, i}, w_{r_2, j}\} x_{r_1, i} x_{r_2, j}$
- Constraints are taken into account in the QUBO formulation to minimize
 - $q = C + \lambda_u C_u + \lambda_t C_t$
 - With :
 - ❖ $C = -\sum_r \sum_{i \in I_r} w_{r,i} x_{r,i}$ is the objective function in the original problem
 - ❖ λ_u, λ_t are penalty weights
- Choice of the penalty weights
 - Sufficiently large enough such that $\hat{x} = \arg \min_x q(x)$ verifies our constraints, i.e. $C_u(\hat{x}) = 0$ and $C_t(\hat{x}) = 0$
 - We can demonstrate that any choice of penalty weight values such that both $\lambda_u > 1$ and $\lambda_t > 1$ gives valid solutions

Mapping a logical QUBO into a physical QUBO

Embedding

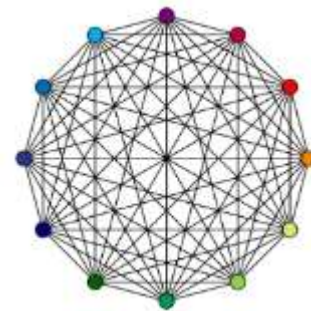
- Due to D-Wave architecture (chimera graph), a physical qubit is not connected to every other qubit
- Embedding is the process of linking physical qubits together to virtually enhance connectivity
- In our case, problem instances need to stay below 80 logical qubits to be embeddable on the D-Wave machine

Weight Distribution

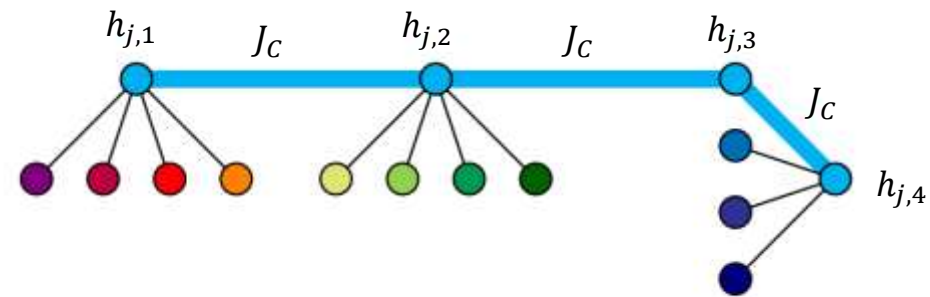
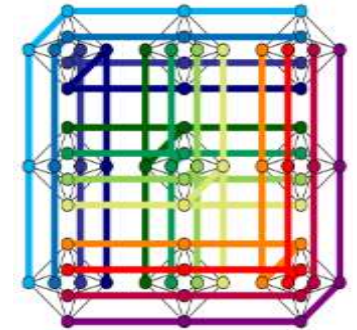
- Couple physical qubits to chain
- Find chain coupling J_C
- Distribute weight h_j
- Classical Approach:
 - Choose J_C according to maximum coupling
 - Split weight equally $h_{j,i} \rightarrow \frac{h_j}{n}$
- Advanced Approaches:
 - Find minimal J_C without breaking chain
 - Map to problem of graph expansion

$$x_i \rightarrow \frac{s_i + 1}{2}$$

$$C = \sum_i h_i s_i + \sum_{ij} J_{ij} s_i s_j \xrightarrow{\text{Embedding}} C' = \sum_a h'_a s'_a + \sum_{ab} J'_{ab} s'_a s'_b$$



Embedding →



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Mission Planning Simulation

Mission Planning problem instances

- Generated thanks to Airbus DS Mission Simulator (TEAM)
- Reduced instances with a small number of requests and a coarse access discretization compared to real operations
- Different scenarios are considered to generate multiple instances, enabling sensitive analysis and statistics on average performance
- Main parameters
 - Number of acquisition requests } → Drives the number of decision variables
 - Access discretization step } → Drives the “NP-hardness” of the planning problem
 - Latitude range for Area of Interest → Drives the “NP-hardness” of the planning problem

Problem Instances

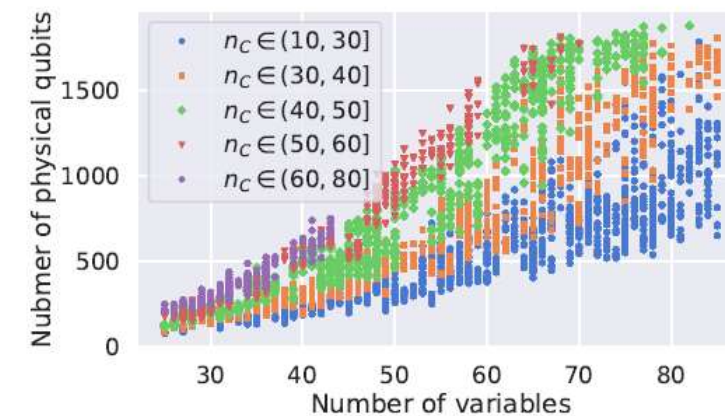
Airbus Amber



Nb of requests = 11
Discretization step = 12s
Latitude range = 10°



Nb of requests = 12
Discretization step = 16s
Latitude range = 1°



Outcome of Embedding

Evaluation on Classical Hardware

Two classical algorithm have been considered

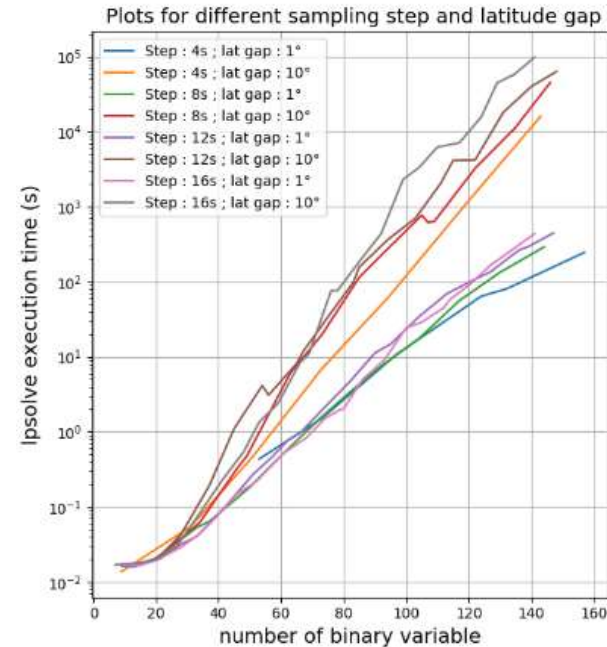
- An exact MIP solver with two variants
 - pairwise exact solver: based on ILP where constraints correspond to pairs of conflicting attempts
 - clique exact solver: based on ILP where constraints are reformulated through the enumeration of all maximal cliques
- A greedy algorithm (similar to operational software), showing a linear runtime (at least for small instances)



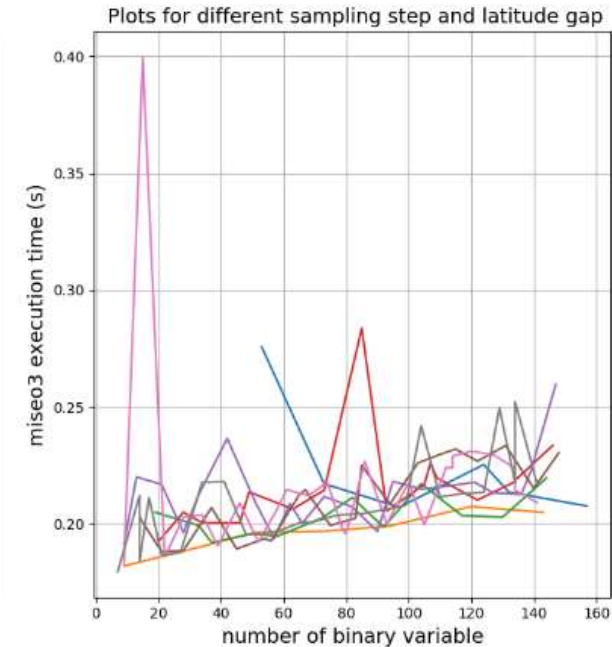
Exact solver



Greedy algorithm



Exact solver
Run-time



Greedy algorithm
Run-time

Evaluation on D-Wave 2000Q Quantum Annealer

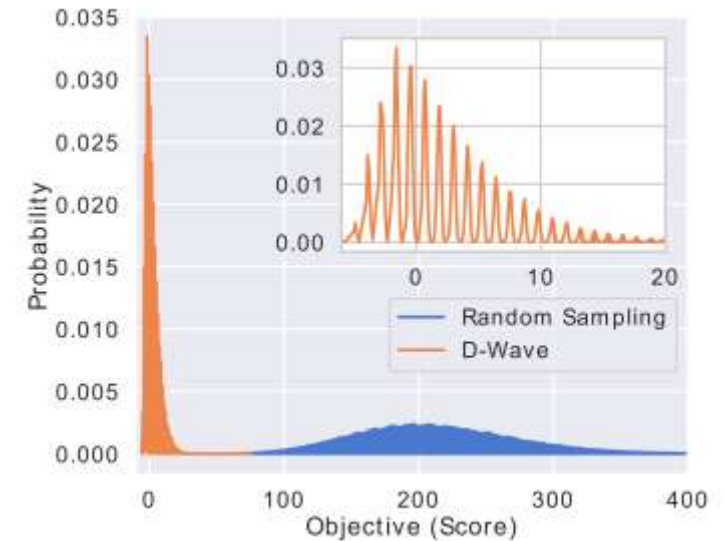
Performance Assessment methodology

- A number of annealing runs is configured
- A success probability is derived: $p = \frac{\# \text{ exact solutions}}{\# \text{ annealing runs}}$
(probability to yield an optimal solution)
- Assuming independence between runs, a time-to-solution with 99% chance of optimality can be expressed

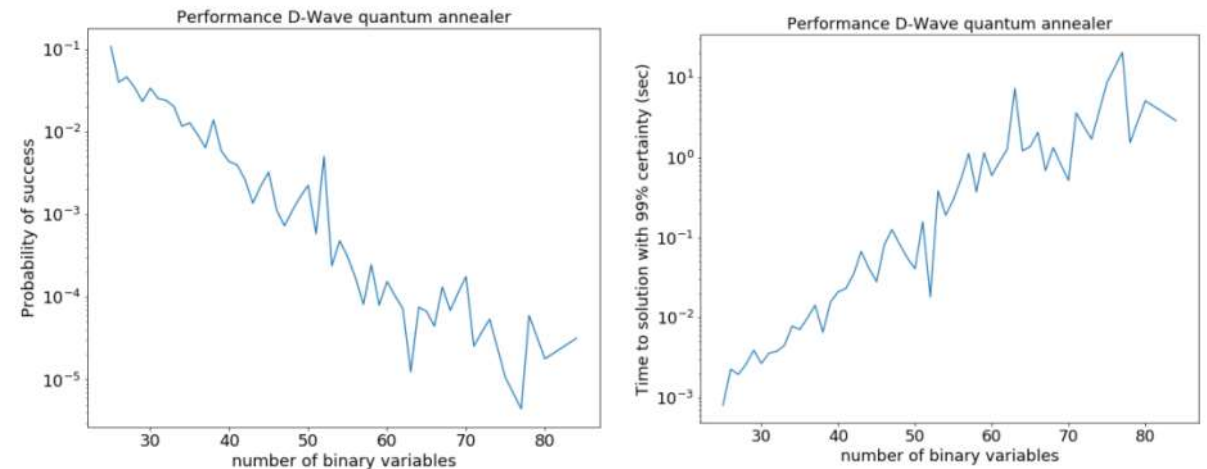
$$\text{as } T_{99} = \frac{\ln(1-0,99)}{\ln(1-p)} T_{\text{Annealing}}$$

D-Wave Configuration

- Number of annealing runs (10000)
- Annealing time (20 μs)
- Choice of intra-logical qubit coupling J_F
- Embeddings: using all 5 D-Wave heuristic embeddings
- Unembedding strategy: majority vote



Random sampler vs D-Wave machine for QUBO resolution



Probability of success

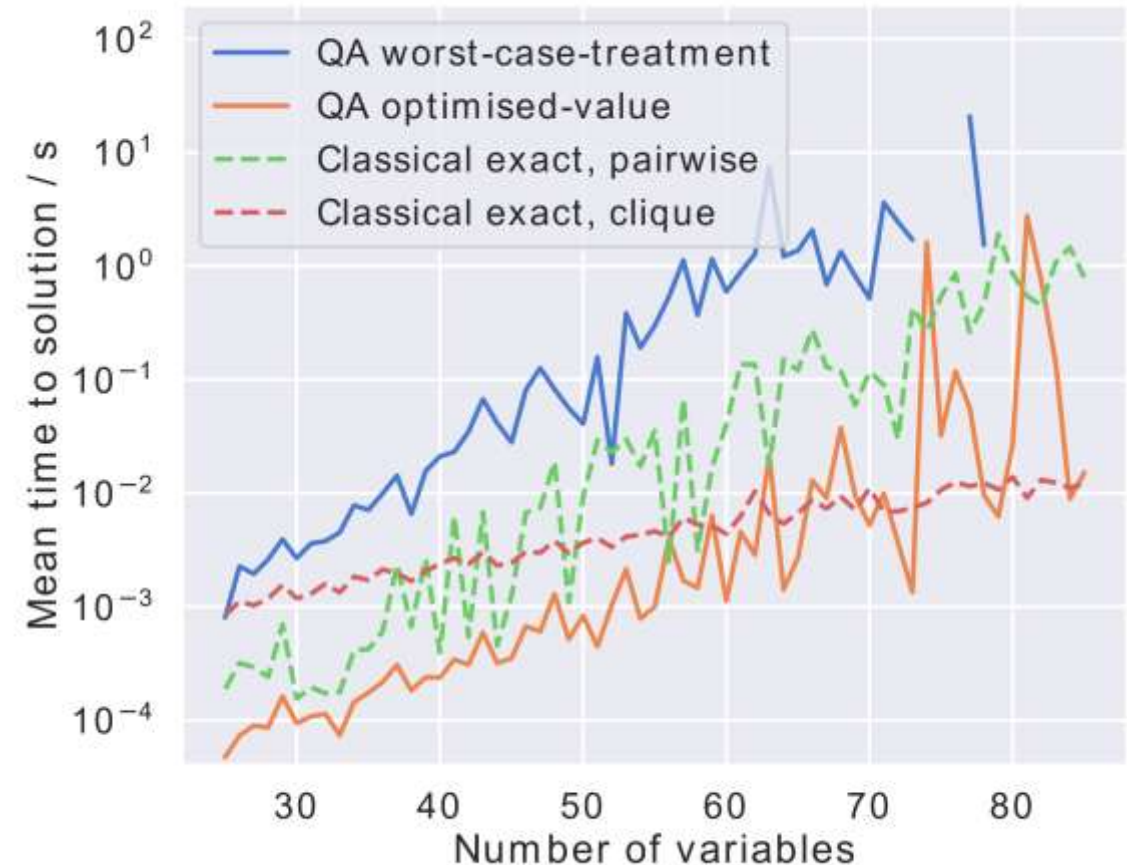


Time to solution with 99%

Benchmark 1: Classical vs Quantum Time to Exact Solution

Time to Exact Solution Benchmark

- Run time is averaged over all problem instances having the same number of binary variables
- Execution time for the pair-wise exact solver increases exponentially with the number of binary variables
- Quantum annealing results (worst-case treatment, i.e. classical weighting approach) shows a similar slope and a constant offset of about one order of magnitude.
- By optimizing the coupling chain strength (optimized-value), quantum annealing performs much better
- Clique exact solver performs better than all other methods for larger instances

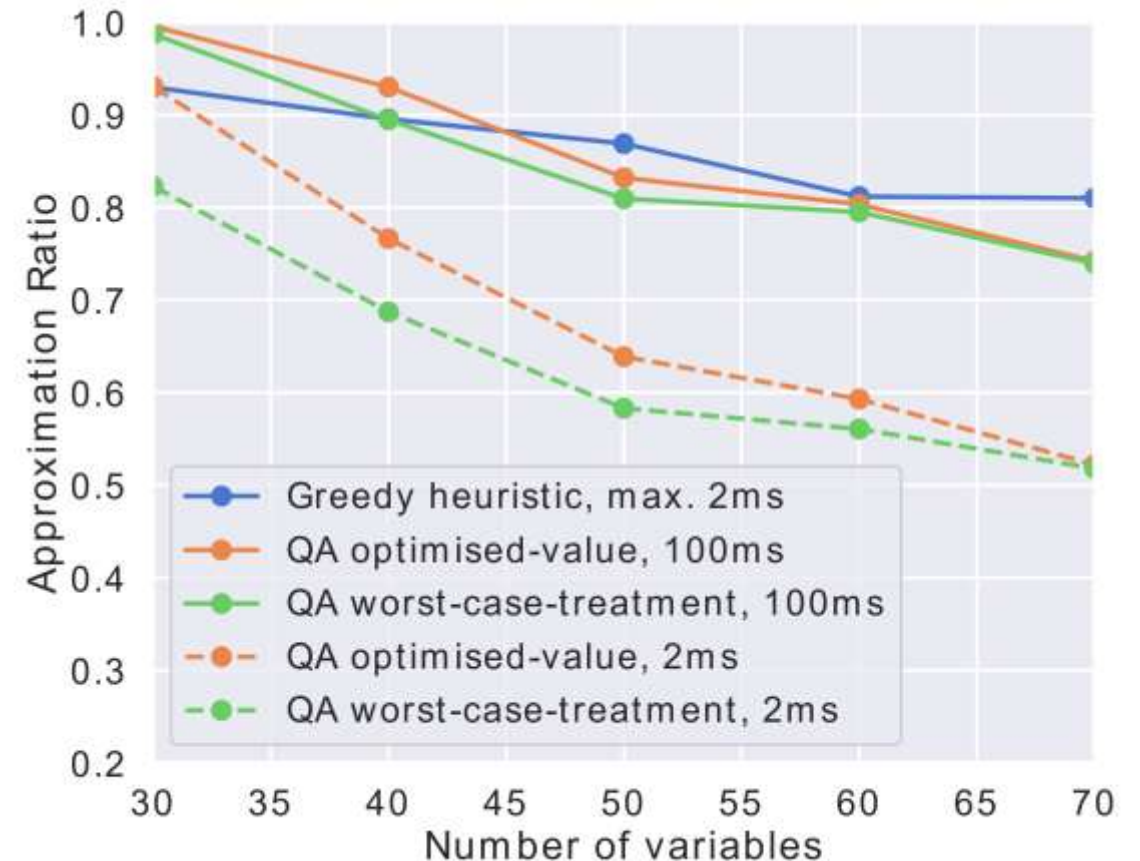


Time to exact solution benchmark

Benchmark 2: Classical vs Quantum Quality of Solution

Quality of Solution Benchmark

- A fixed time budget is allocated to the solver (.i.e a fixed number of runs for the quantum annealer)
- The approximation ratio corresponds to the objective value of the best found solution divided by the optimal objective value
- The greedy heuristic outperforms the quantum annealer for similar execution times
- Only for larger execution times, the quantum annealer yields better results than the greedy heuristic for smaller instances



Quality of solution benchmark

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Conclusion and Perspectives

arXiv paper: <https://arxiv.org/abs/2006.09724>

Technical achievements

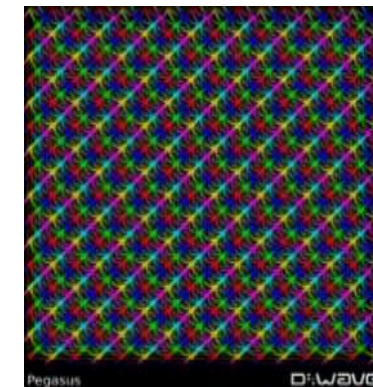
- Classical vs Quantum benchmarks for a broad range of (small) satellite mission planning problems
- Limited qubit connectivity, **precision issues** and coherence time remain a major bottleneck for the D-Wave 2000Q processor.
- Although **no quantum speedup** was observed, the run-time performance on D-Wave Q Annealer (at its current scale) is very promising

Perspectives

- Extra research will be required to make a better use of Q technology (embedding techniques and mitigation of precision/errors for QA)
- To draw further conclusions, we need the Q technology (HW and SW) to **increase in maturity**, which will happen in a short timeframe
 - D-Wave Pegasus showcasing 5000 qubits and 16-connectivity
 - QAOA on Google Sycamore and IBM Q 53-qubit machines



ESA Phi-Week



D-Wave Pegasus



Google Sycamore

Thank you



DLR

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